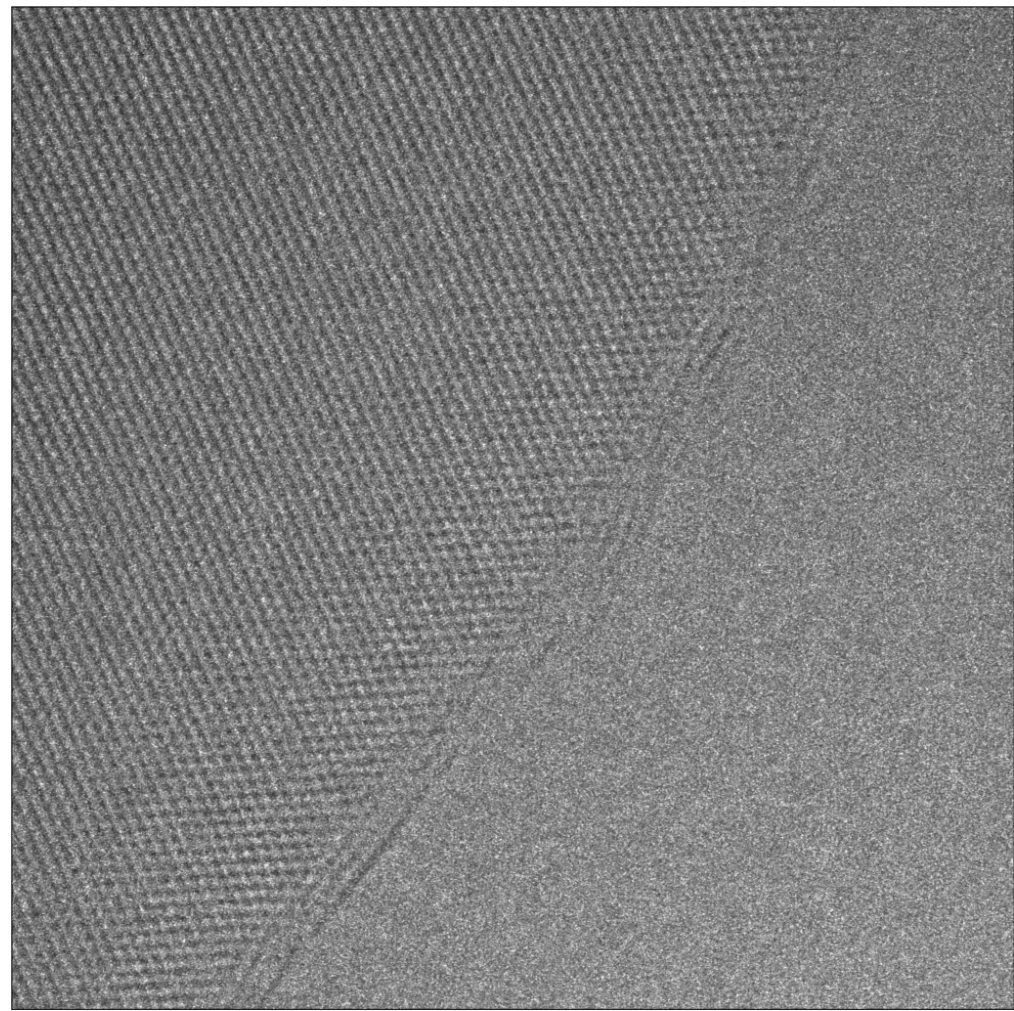


Parallel Prefix Algorithms for the Registration of Arbitrarily Long Electron Micrograph Series

Marcin Copik, Paolo Bientinesi (advisor), Benjamin Berkels (advisor)

Motivation



The process of aluminum oxidization, acquired with a transmission electron microscope. Courtesy of the University of Manchester.

- ▶ The interaction of electrons with the specimen is mapped to a display device
- ▶ The quality of images is limited by changes induced by electrons and movement of the sample
- ▶ Solve both problems by acquiring a series of low dose frames and processing it to extract information
- ▶ Image registration is computationally expensive - can we parallelize that process?

Image registration

Goal: find rigid deformations $\phi_{0,i}$ for frames $f_1, \dots, f_n$

$$f_0 \approx f_i \circ \phi_{0,i}$$

Neighboring frames: apply registration function **A**

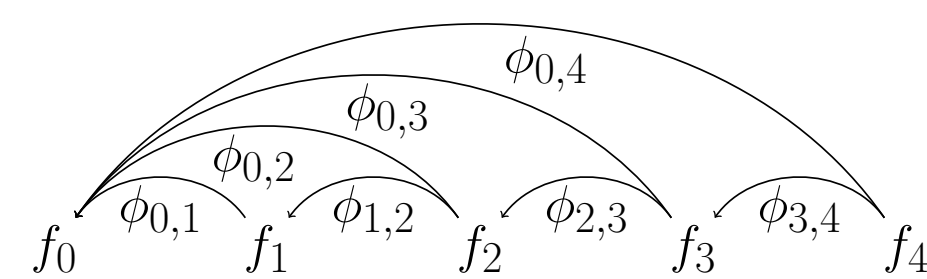
$$\phi_{i,i+1} = \mathbf{A}(f_i, f_{i+1})$$

Non-neighboring frames: use registration function **B**

$$\phi_{i,j} = \mathbf{B}(\phi_{i,k}, \phi_{k,j})$$

Prefix sum registration: with a binary operator $\odot_B$

$$\phi_{0,i} = \phi_{0,1} \odot_B \phi_{1,2} \odot_B \dots \odot_B \phi_{i-1,i}$$



Parallel image registration

- ▶ A **trivially parallelizable** preprocessing step on neighboring pairs of frames

$$f_0, f_1, \dots, f_n \rightarrow \phi_{0,1}, \phi_{1,2}, \dots, \phi_{n-1,n}$$

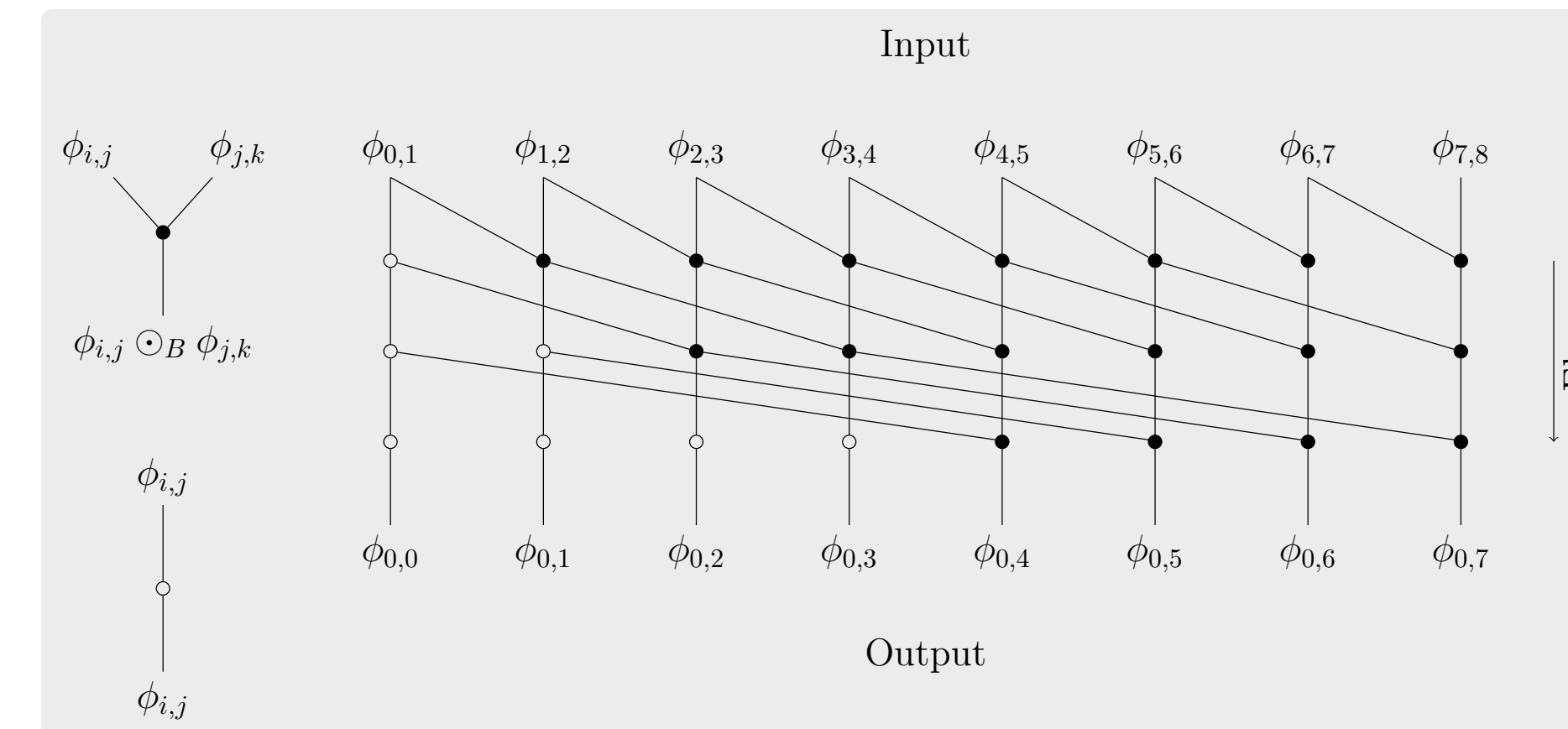
- ▶ A prefix sum registers each frame to the first image, producing a new sequence

$$\phi_{0,1}, \phi_{0,2}, \dots, \phi_{0,n}$$

a parallel prefix sum is necessary

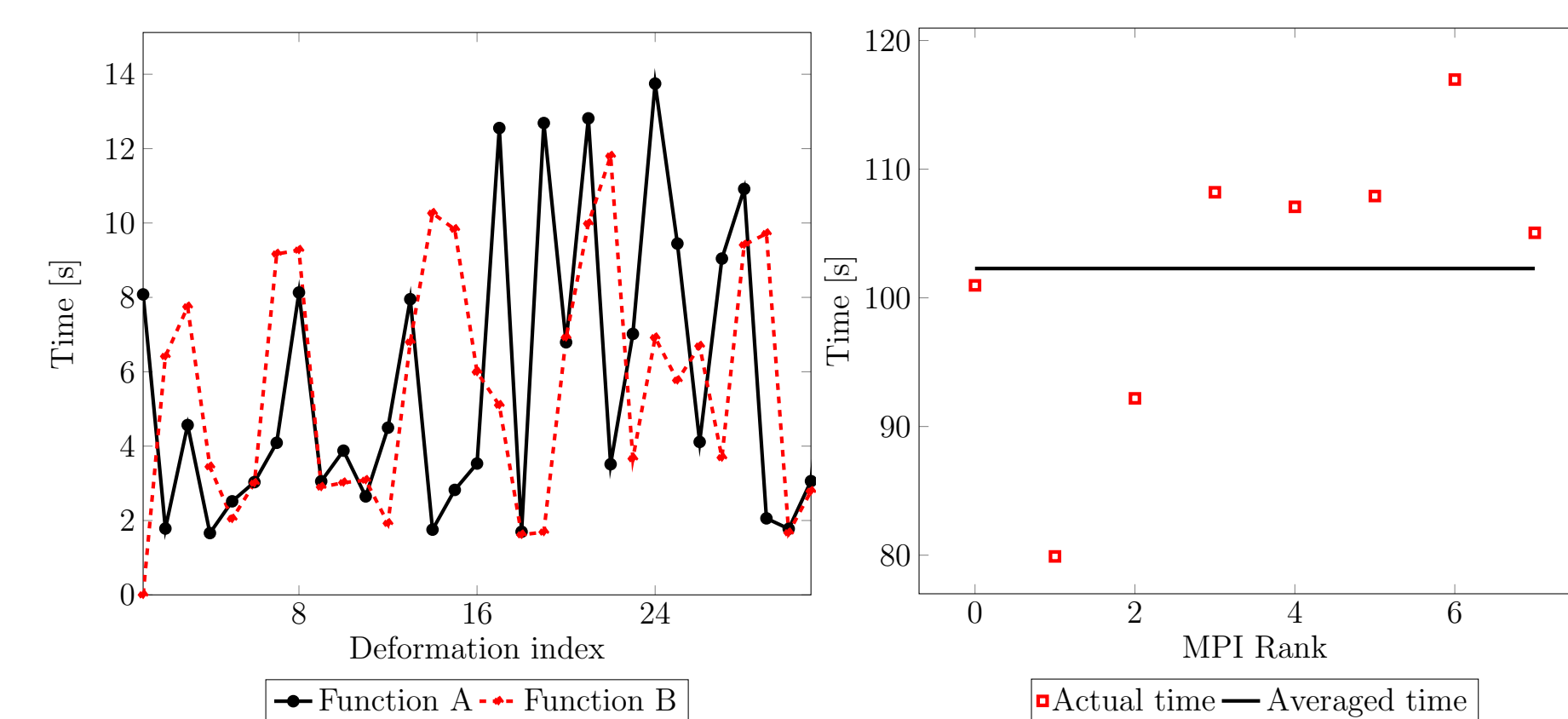
Parallel prefix sum

- ▶ Multiple algorithms for of a parallel prefix sum exist
- ▶ A decrease in span is achieved by performing more work
- ▶ Scan algorithms with $\log_2 N$ span are optimal, but they can not be work-efficient



Kogge-Stone(Hills-Steele) parallel prefix sum
 $S(N) = \log_2 N$

Image registration operator



(1) Registration times on 64 images (2) Function B on 128 deformations on 8 MPI ranks

Existing parallel prefix sum algorithms

- ▶ Are simple and not computationally intensive
- ▶ Have a stable and deterministic execution time

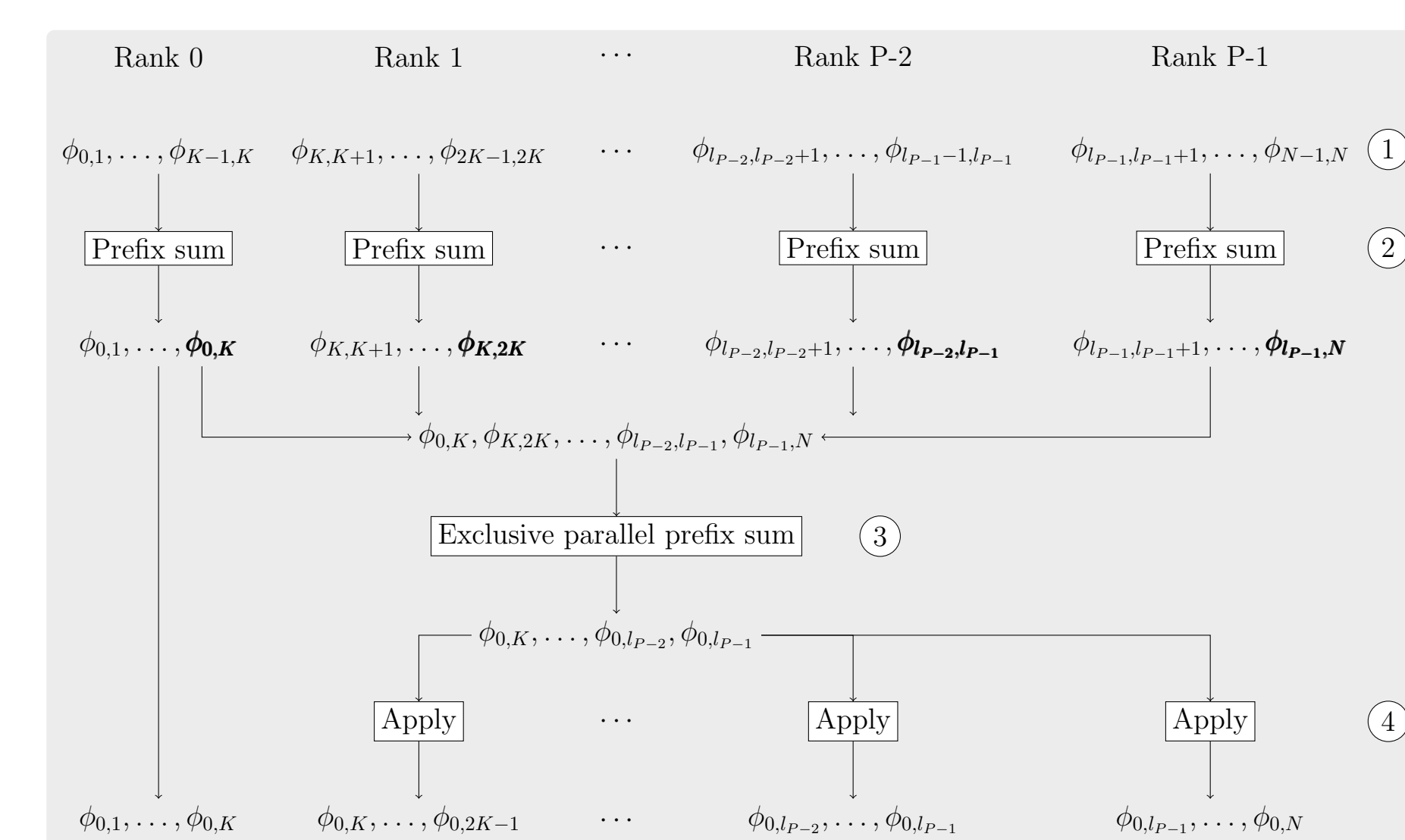
In our problem, the image registration operator

- ▶ Is complex and computationally intensive (1)
- ▶ Has an unpredictable and unbalanced execution time (2)

Parallel prefix sum for image registration

- ▶ Communication cost $\ll$ computation cost
 $\Rightarrow$ do not optimize for communication
- ▶ Reduce computation time
 $\Rightarrow$ minimize span
- ▶ Variations in execution time are not known a priori
 $\Rightarrow$ distribute data equally between MPI ranks

Distributed prefix sum



- Rank J obtains a sequence of data $\phi_{l_j, l_j+1}, \phi_{l_j+1, l_j+2}, \dots, \phi_{l_j+1, l_j+1}, l_j+1$ with $K = \frac{N}{P}$ elements, where $l_j = J \cdot K$ is a left index for rank J
- Rank J performs a local prefix sum, reducing sequence to ϕ_{l_j, l_j+1}

$$S(N, P) = \frac{N}{P} - 1$$

- A global exclusive prefix sum on reduced values ϕ_{l_j, l_j+1} is performed, rank J obtains a result ϕ_{0, l_j}

$$S(N, P) = C_1 \log_2 P$$

- The global result ϕ_{0, l_j} and local values ϕ_{l_j, l_j+i} are combined to form final results ϕ_{0, l_j+i}

$$S(N, P) = \frac{N}{P}$$

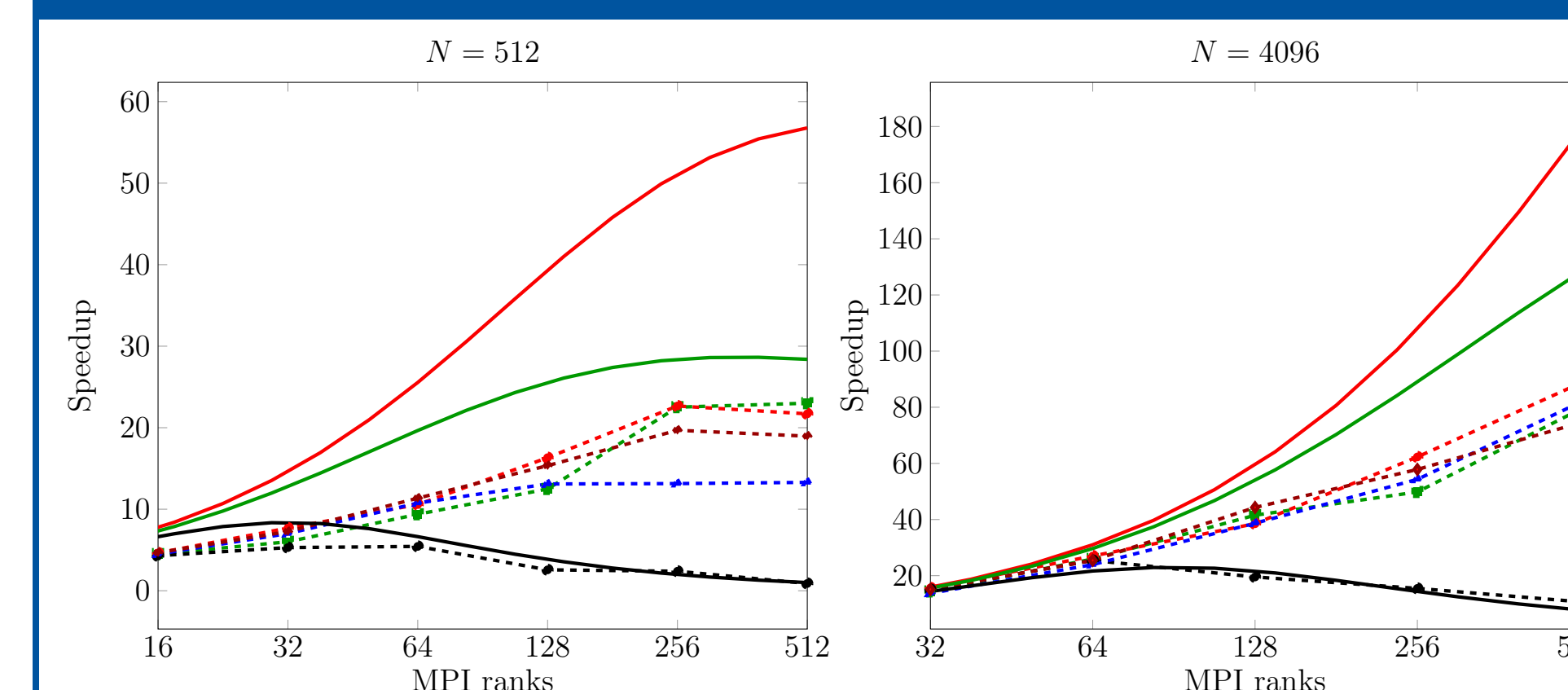
The span of a prefix sum for N images on P ranks

$$S(N, P) = \frac{N}{P} + C_1 \log_2 P + \frac{N}{P} + C_2$$

The attainable speedup is bounded from above by

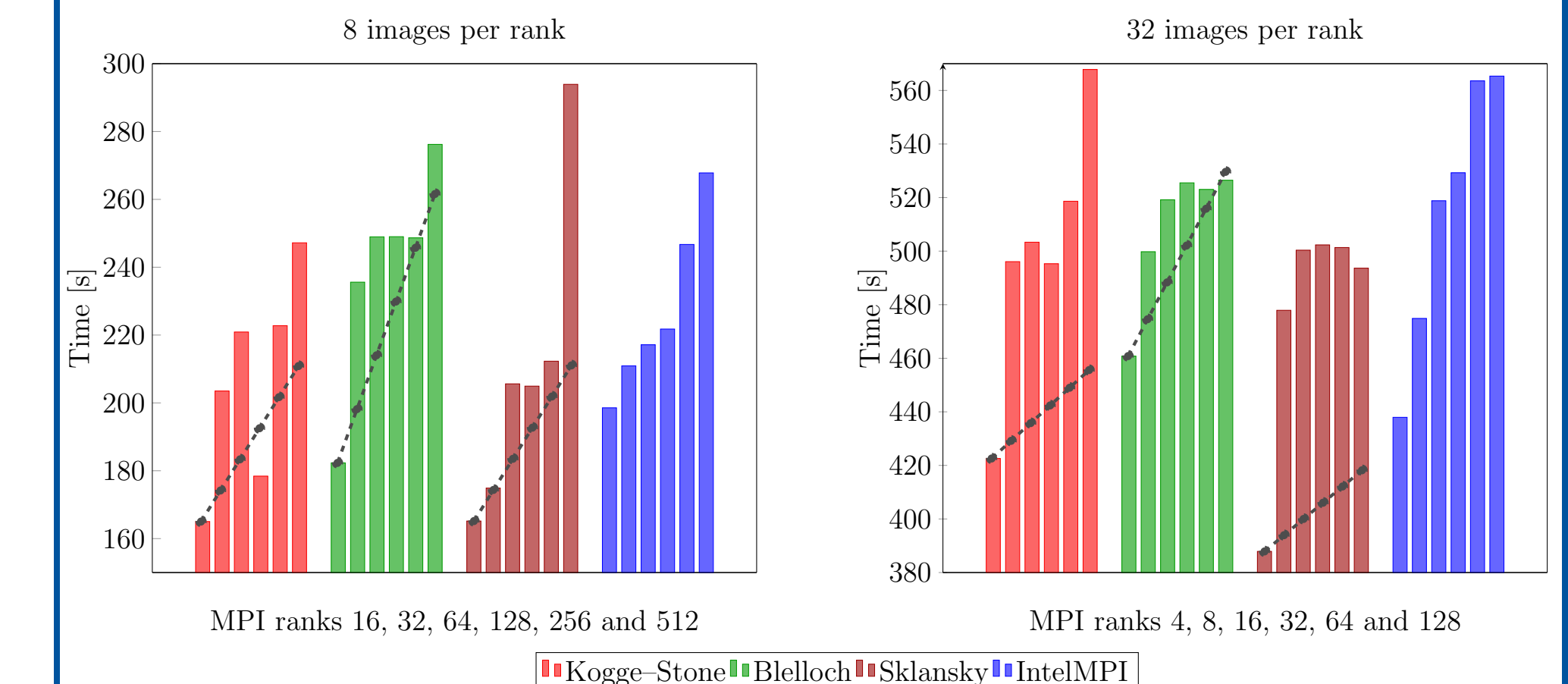
$$SP(N, P) = \frac{S(N, 1)}{S(N, P)} = \frac{N - 1}{\frac{2N}{P} + C_1 \log_2 P + C_2}$$

Strong scaling



Theoretical and measured speedups are represented with solid and dashed lines, respectively.

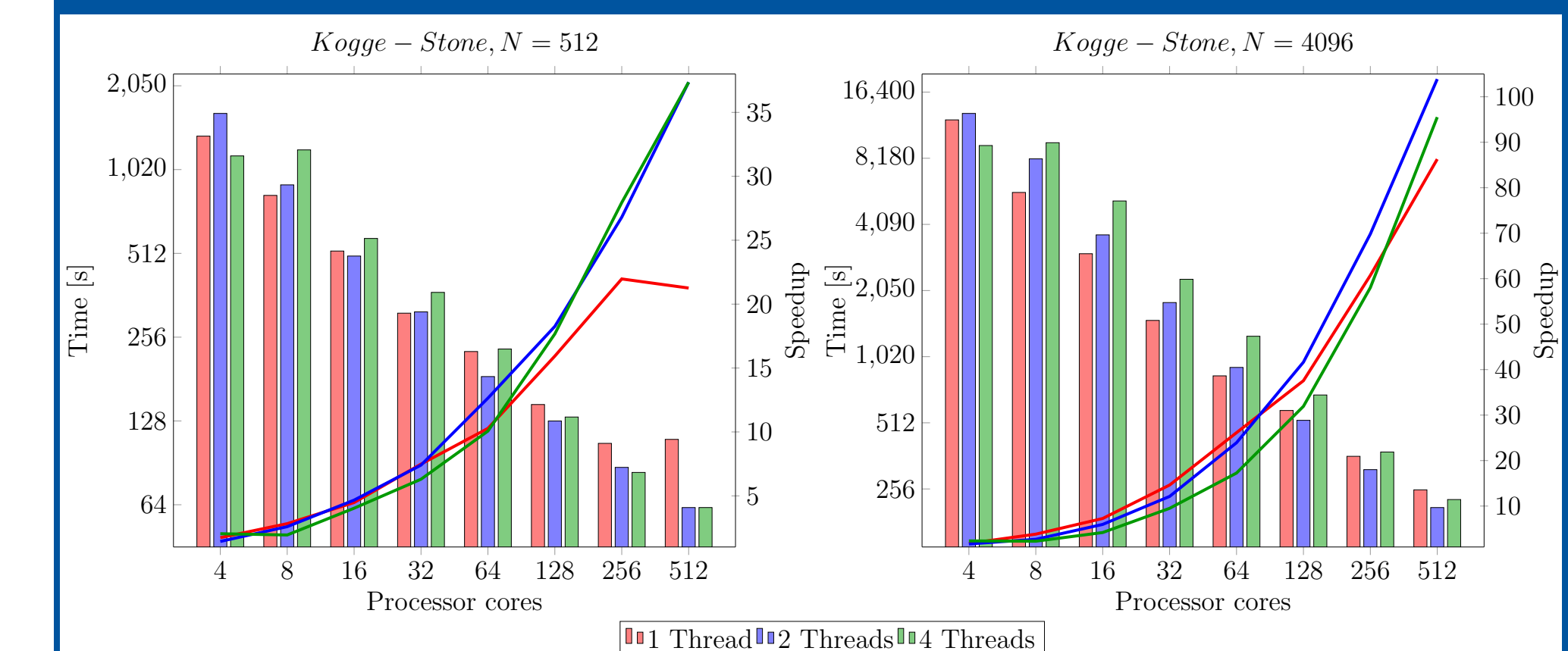
Weak scaling



Linear weak scaling is not attainable, since

$$S(k \cdot N, k \cdot P) = S(N, P) + \log_2 k \cdot C_1$$

Multithreaded registration operator



Allocate $\frac{P}{T}$ ranks with T threads per each MPI process.

Summary

- ▶ The benchmarking has been performed on Intel Xeon E5-2680 IvyBridge processors with GCC 5.3 compiler
- ▶ Measured performance does not meet theoretical predictions because of the unbalanced workload
- ▶ Dynamic scheduling is limited by the nature of prefix sum and the distributed environment
- ▶ The performance is improved by shifting resources to multithreaded operator parallelization
- ▶ Custom scan implementations outperform MPI_Scan and MPI_Exscan in OpenMPI 1.10, IntelMPI 2017.1
- ▶ The image registration procedure is a novel example of parallel prefix sum with a non-trivial sum operator

References

- ▶ Benjamin Berkels et al. (2014). "Optimized imaging using non-rigid registration". In: *Ultramicroscopy* 138, pp. 46–56
- ▶ G. E. Blelloch (1989). "Scans As Primitive Parallel Operations". In: *IEEE Trans. Comput.* 38.11, pp. 1526–1538



Marcin Copik
E-Mail: mcopik@gmail.com

High Performance and Automatic Computing
Schinkelstr. 2, 52062 Aachen, GERMANY
www.hpac.rwth-aachen.de
Telefon: +49 241 80 99 131, Fax: +49 241 80 628498